

# CONTRASTIVE EVIDENCE IN THE AGM FRAMEWORK

JAKE CHANDLER  
University of Leuven

## 1. Preliminaries

$\mathcal{L}$ : arbitrary propositional ('factual') language, constructed by means of the standard Boolean connectives  $\{\vee, \wedge, \neg, \rightarrow\}$

$\mathcal{L}_E$ : smallest extension of  $\mathcal{L}$  that includes all sentences of the form  $\varphi \triangleright \psi$  and  $\varphi \triangleright_\chi \psi$  (where  $\varphi, \psi, \chi \in \mathcal{L}$ )

$K$ : arbitrary belief set, a subset of  $\mathcal{L}_E$

$\mathbb{K}$ : set of all rationally permissible belief sets

$M$ : arbitrary belief revision model, a pair of functions  $\langle *, \dot{-} \rangle$ , such that  $*$  (the revision function) and  $\dot{-}$  (the contraction function) map elements of  $\mathbb{K} \times \mathcal{L}$  onto  $\mathbb{K}$

$\mathbb{M}$ : set of all rationally permissible belief revision models

$\text{Cn}$ : consequence operator, mapping subsets of  $\mathcal{L}_E$  onto other such subsets

## 2. Principles

Consequence:

(DT) For all  $A, B \in \mathcal{L}$  and  $\Gamma \subseteq \mathcal{L}_E$ , if  $B \in \text{Cn}(\Gamma \cup \{A\})$ , then  $A \rightarrow B \in \text{Cn}(\Gamma)$

**(SUP)** For all  $\Gamma \subseteq \mathcal{L}_E$ ,  $\text{Cn}_0(\Gamma_0) \subseteq \text{Cn}(\Gamma)$ , where  $\Gamma_0$  denotes the factual subset of  $\Gamma$  and  $\text{Cn}_0$  returns the set of its classical consequences

Belief sets:

**(CL)** For all  $K \in \mathbb{K}$ ,  $\text{Cn}(K) \subseteq K$

**(CON)** For all  $K \in \mathbb{K}$  and  $A \in \mathcal{L}$ , there is no finite  $\Gamma \subseteq K$  such that  $A \wedge \neg A \in \text{Cn}(\Gamma)$ .

Revision & contraction:<sup>1</sup>

**(AGM\*2)** If  $\neg A \notin \text{Cn}(\emptyset)$ , then  $A \in K * A$ .

**(AGM\*3)** If  $B \in K * A$ , then  $B \in \text{Cn}(K \cup \{A\})$

**(AGM\*V)** If  $\neg A \in \text{Cn}(\emptyset)$  then  $K * A = K$ .

**(AGM\*6)** If  $\text{Cn}(A) = \text{Cn}(B)$ , then  $K * A = K * B$

**(AGM\*7)**  $K * (A \wedge B) \subseteq \text{Cn}(K * A \cup \{B\})$

**(CM)** If  $B \in K * A$  and  $C \in K * A$ , then  $C \in K * A \wedge B$

**(I\*1)** If  $C \in \text{Cn}(A)$ , then  $B \in K * A$  iff  $B \in (K * C) * A$

**(I\*3)** If  $B \in K * A$ , then  $B \in (K * B) * A$

**(AGM  $\dot{-}$ 2)** If  $B \in K \dot{-} A$ , then  $B \in K$

**(AGM  $\dot{-}$ 3)** If  $A \notin K$  or  $A \in \text{Cn}(\emptyset)$ , then  $K \dot{-} A = K$

**(AGM  $\dot{-}$ 4)** If  $A \notin \text{Cn}(\emptyset)$ , then  $A \notin K \dot{-} A$

**(AGM  $\dot{-}$ 5)** If  $\text{Cn}(A) = \text{Cn}(B)$ , then  $K \dot{-} A = K \dot{-} B$

**(AGM  $\dot{-}$ 6)** If  $B \in K$ , then  $B \in \text{Cn}((K \dot{-} A) \cup \{A\})$

**(I  $\dot{-}$ 2)** If  $C \in \text{Cn}(A)$ , then  $B \in K * A$  iff  $B \in (K \dot{-} C) * A$

**(I  $\dot{-}$ 4)** If  $B \notin K * A$ , then  $B \notin (K \dot{-} B) * A$

**(I  $\dot{-}$ 4+)** If  $B \in \text{Cn}(C)$ , then if  $B \notin (K \dot{-} C) * A$ , then  $B \notin (K \dot{-} B) * A^2$

<sup>1</sup>All principles are to be read as holding for all  $K \in \mathbb{K}$ ,  $A, B, C \in \mathcal{L}$  and  $*$ ,  $\dot{-}$  that are members of some  $M \in \mathbb{M}$ .

<sup>2</sup>Entails (I  $\dot{-}$ 4).

Evidential beliefs:

**(BEV)**  $A \triangleright B \in K$  iff  $B \in (K \dot{-} B) * A$

**(TEV)** Where  $\neg C \in \text{Cn}(B)$ ,  $A \triangleright_C B \in K$  iff  $B \in ((K \dot{-} \neg C) * B \vee C) * A$

**(TR1)** If  $C \in \text{Cn}(B)$ , then  $A \triangleright C \in \text{Cn}(A \triangleright B)$ .

**(TR1')** If  $C \in \text{Cn}(B)$  and  $A \triangleright B \in K$ , then  $A \triangleright C \in K$ .

### 3. Observations and selected proofs

**Obs1:** Given (CL), (CON), (DT), (AGM\*2), (AGM\*3), (AGM\*6), (AGM\*7), (AGM $\dot{-}$ 3), (AGM $\dot{-}$ 5) and (I $\dot{-}$ 4+), (BEV) entails:

**(ABS)** If  $A \triangleright B \in K$ , then  $A \triangleright (A \wedge B) \in K$

**(CA)**  $A \triangleright C \in K$  and  $B \triangleright C \in K$ , then  $(A \vee B) \triangleright C \in K$

**(CC)** If  $A \triangleright B \in K$  and  $A \triangleright C \in K$ , then  $A \triangleright (B \wedge C) \in K$

**(CN)** If  $\neg A \notin \text{Cn}(\emptyset)$  and  $B \in \text{Cn}(\emptyset)$ , then  $A \triangleright B \in K$

**(ID)** If  $A \notin \text{Cn}(\emptyset)$ , then  $A \triangleright A \in K$

**(LLE)** If  $(A \leftrightarrow C) \in \text{Cn}(\emptyset)$ , then  $A \triangleright B \in K$  iff  $C \triangleright B \in K$ .

**(MP)** If  $A \triangleright B \in K$  then  $A \rightarrow B \in K$

**(NEC)** If  $\neg B \in \text{Cn}(\emptyset)$  then  $A \triangleright B \notin K$

**(RLE)** If  $(B \leftrightarrow C) \in \text{Cn}(\emptyset)$ , then  $A \triangleright B \in K$  iff  $A \triangleright C \in K$

**(RRT)** If  $B \notin K$ , then  $A \triangleright B \in K$  iff  $B \in K * A$

**(SCC)** If  $A \triangleright B \in K$  and  $\neg B \vee \neg C \in \text{Cn}(\emptyset)$ , then  $A \triangleright C \notin K$

**(SUP $\triangleright$ )** If  $A \notin \text{Cn}(\emptyset)$  and  $B \in \text{Cn}(A)$ , then  $A \triangleright B \in K$

*Proof:* Available on request.

**Obs2:** Given (CL), (CON), (AGM\*2), (AGM\*3), (AGM $\dot{-}$ 4) and (I $\dot{-}$ 4+), (BEV) is inconsistent, on pains of various triviality results, with:

**(CCC)** If  $A \triangleright B \in K$  and  $B \in \text{Cn}(C)$ , then  $A \triangleright C \in K$

(CS) If  $A, B \in K$ , then  $A \triangleright B \in K$

(EEM) Either  $A \triangleright B \in K$  or  $A \triangleright \neg B \in K$

(SYM) If  $A \triangleright B \in K$ , then  $B \triangleright A \in K$

*Proof:* Available on request.

**Obs3:** (BEV), in the presence of (TR1'), (CL), (SUP), (AGM\*2), and (AGM\*4) entails:

(TRIV) There is no  $K \in \mathbb{K}$ , and  $A, B, C \in \mathcal{L}$  such that  $A \in K$ ,  $B, \neg B \notin K$  but  $B \triangleright A \notin K$ .

*Proof:* See Chandler (2012), fn. 27.

**Obs4:** Given (AGM \* 2), (AGM \* 3), (AGM \* V), (CM), (I\*1), (AGM  $\dot{\vdash}$  2), (AGM  $\dot{\vdash}$  3), (AGM  $\dot{\vdash}$  4), (AGM  $\dot{\vdash}$  6), (I  $\dot{\vdash}$  2), (I  $\dot{\vdash}$  4), (CL), (DT) and (SUP), (BEV) entails

(INT1) If  $C \in \text{Cn}(B)$  and  $A \triangleright B \in K$ , then  $A \triangleright C \notin K$  iff  $A \triangleright B \notin K \dot{\vdash} C$ .

*Proof:* See Chandler (2012), proof of Observation 5.

**Obs5:** Given (AGM\*V), (BEV) and (TEV) jointly entail that  $A \triangleright B \in K$  iff  $A \triangleright_{\neg B} B \in K$ .

*Proof:* By (SUP),  $\text{Cn}(B) = \text{Cn}(\neg\neg B)$ , so by (AGM  $\dot{\vdash}$  6)  $(K \dot{\vdash} B) = (K \dot{\vdash} \neg\neg B)$ . Furthermore, by (SUP) again,  $B \vee \neg B \in \text{Cn}(\emptyset)$ , so by (AGM\*V),  $(K \dot{\vdash} \neg\neg B) = (K \dot{\vdash} \neg\neg B) * B \vee \neg B$ . Putting this together,  $(K \dot{\vdash} B) = (K \dot{\vdash} \neg\neg B) * B \vee \neg B$ , and hence  $(K \dot{\vdash} B) * A = ((K \dot{\vdash} \neg\neg B) * B \vee \neg B) * A$ . By (BEV) and (TEV), this is equivalent to  $A \triangleright B \in K$  iff  $A \triangleright_{\neg B} B \in K$ . QED.

**Obs6:** Given (DT), (CL), (I\*1), (I\*3), (AGM\*7), (CM), (I  $\dot{\vdash}$  2) and (I  $\dot{\vdash}$  4), (TEV) entails:

(INT2) If  $C \in \text{Cn}(B)$ ,  $A \triangleright B \in K$ , then (1)  $A \triangleright B \notin K \dot{\vdash} C$  iff (2)  $A \triangleright_{\neg C} B \notin K$ .

*Proof:* We first note that the following principle, which we shall occasionally use later on, is a well-known weakening of (AGM\*7):

(CUT) If  $B \in K * A$ , then  $K * (A \wedge B) \subseteq K * A$

With this in hand, we now assume  $C \in \text{Cn}(B)$ . We first establish the following lemma:

*Lemma 1:* If (i)  $B \in K * A$ , then (ii)  $K * A = (K * C) * A$

Assume (i) for conditional proof. By (i), the fact that  $C \in \text{Cn}(B)$  and (CL), it follows that

$$C \in K * A \quad (1)$$

From (1), it then follows by (CM) and (CUT) that

$$K * A = K * A \wedge C \quad (2)$$

From (1), but this time by (DP\*3), we obtain

$$C \in (K * C) * A \quad (3)$$

And, from (3), it then follows by (CM) and (CUT) that

$$(K * C) * A \wedge C = (K * C) * A \quad (4)$$

Finally, since  $C \in \text{Cn}(A \wedge C)$ , it follows from (DP\*1) that

$$K * A \wedge C = (K * C) * A \wedge C \quad (5)$$

Putting together (2), (4) and (5), we recover the desired result that (ii)  $K * A = (K * C) * A$ . This completes the proof of Lemma 1.

We now derive a second lemma:

*Lemma 2:* If (i)  $A \triangleright B \in K$ , then (ii)  $B \in (K \dot{-} C) * A \wedge C^3$

Assume (i) for conditional proof. It follows, by the left-to-right direction of (BEV), that

$$B \in (K \dot{-} B) * A \quad (6)$$

From (6), and the fact that  $C \in \text{Cn}(B)$ , it follows by (CL) that

$$C \in (K \dot{-} B) * A \quad (7)$$

From (6), (7) and (CM), it then follows that

$$B \in (K \dot{-} B) * A \wedge C \quad (8)$$

---

<sup>3</sup>This is Lemma 1 of the proof of Observation 5 in Chandler (2012).

Now assume for reductio that  $B \notin (K \dot{-} C) * A \wedge C$ . Since  $C \in \text{Cn}(A \wedge C)$ , it follows by (I  $\dot{-}$ 2) that

$$B \notin K * A \wedge C \quad (9)$$

By (9) and (I  $\dot{-}$ 4), it then follows that

$$B \notin (K \dot{-} B) * A \wedge C \quad (10)$$

Contradiction. Hence (ii):  $B \in (K \dot{-} C) * A \wedge C$ . This completes the proof of Lemma 2.

Finally, we derive a third and final lemma:

*Lemma 3:* If  $B \in (K \dot{-} C) * A \wedge C$ , then  $A \triangleright B \notin K \dot{-} C$  iff  $A \triangleright_{\neg C} B \notin K$

Assume  $B \in (K \dot{-} C) * A \wedge C$  for conditional proof. By (AGM\*7)

$$B \in \text{Cn}((K \dot{-} C) * A \cup \{C\}) \quad (11)$$

It follows from (11), by the deduction theorem (DT), which Cn satisfies, that

$$B \vee \neg C \in (K \dot{-} C) * A \quad (12)$$

Obviously,  $B \vee \neg C \in \text{Cn}(B \vee \neg C)$ . So by (12) and Lemma 1, we have  $(K \dot{-} C) * A = ((K \dot{-} C) * B \vee \neg C) * A$  and hence

$$B \in (K \dot{-} C) * A \text{ iff } B \in ((K \dot{-} C) * B \vee \neg C) * A \quad (13)$$

By (BEV) and (TEV), (13) is equivalent to:

$$A \triangleright B \notin K \dot{-} C \text{ iff } A \triangleright_{\neg C} B \notin K \quad (14)$$

This completes the proof of Lemma 3.

Obs6 then follows from the conjunction of Lemmas 2 and 3. QED.

## References

- Achinstein (2001). *The Book of Evidence*. Oxford: OUP.  
 Chandler, J. (2007). Solving the Tacking Problem with Contrast Classes. *British Journal for the Philosophy of Science*, 58(3):489–502.

- Chandler, J. (2009). The Transmission of Support: A Bayesian Re-analysis. *Synthese*, 176(3):333–343.
- Chandler, J. (2010). Contrastive Confirmation: Some Competing Accounts. *Synthese* online first. DOI: 10.1007/s11229-010-9845-9.
- Chandler, J. (2012). Transmission-Failure, AGM-Style. *Erkenntnis*, online first. DOI: 10.1007/s10670-012-9364-9.
- Darwiche, A. & Pearl, J. (1997). On the Logic of Iterated Belief Revision. *Artificial Intelligence* 89: 1–29.
- Fitelson, B. (2007). Likelihoodism, Bayesianism, and Relational Confirmation. *Synthese*, 156(3):473–489.
- Fitelson, B. (2012). Contrastive Bayesianism. In M. Blauw (ed.) *Contrastivism in Philosophy*, London: Routledge.
- Gärdenfors, P. (1984). Epistemic Importance and Minimal Changes of Belief. *Australasian J. Phil.* 62:136–157.
- Okasha, S. (2004). Wright on the Transmission of Support: a Bayesian Analysis. *Analysis*, 64(282):139–146.
- Sober, E. (1994). Contrastive empiricism. In his *From A Biological Point of View: Essays in Evolutionary Philosophy*, pp. 114–135. Cambridge: Cambridge University Press.
- Wright, C. (1985). Facts and Certainty. *Proceedings of the British Academy*, 71:429–472.
- Wright, C. (2002). (Anti-)Sceptics Simple and Subtle: G. E. Moore and John McDowell. *Philosophy and Phenomenological Research*, 65(2):330–348.