

MODELING MODAL AGNOSTICISM

J. Chandler*

*Centre for Logic and Analytic Philosophy
University of Leuven

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Introduction

- Large literature on rational acceptability of *conditionals* wrt epistemic states modeled as orderings of worlds by comparative plausibility.
- Comparatively neglected topic: rational acceptability of *modals*, of the form 'It might/must be the case that P '. (One of Hansson's 'ten philosophical problems in belief revision'; Hansson [2003])
- Received view: Levi [1988], whose acceptance conditions impose strong constraints on rational agents.
- In a recent *Mind* article, Sorensen [2009] puts forward considerations that suggest that these constraints are *too* strong.

Introduction

- In this talk: discussion of impact of Sorensen's view on the standard world-order model of epistemic states.
- I offer a required generalization of the standard model.
- I also briefly discuss an associated modal logic with a clear supervaluationist flavour.

Outline

Formal preliminaries

Modals beliefs and factual beliefs

Epistemic states

Logics

Concluding comments

Languages

- \mathcal{L}_0 Arbitrary ‘factual’ propositional language, generated from a finite (!) set \mathcal{A} of atomic sentences using the Boolean connectives $\{\wedge, \vee, \rightarrow, \neg\}$.
- W_0 Set of valuations of \mathcal{L}_0 .
- $[[\varphi]]$ Set of all $x \in W_0$, such that $x \models \varphi$.
- \mathcal{L}_M Extension of \mathcal{L}_0 , adding a unary possibility connective \diamond .
- Intended interpretation:
 - ‘It might be the case that...’ (\neq ‘It might *have been*...’!!)
 - ‘There is a possibility that...’ (\neq ‘There *would have been*...’!!)
- \square Shorthand for $\neg \diamond \neg$

Consequence

- **Cn** Consequence operator:
 - Function from $\wp(\mathcal{L}_M)$ to $\wp(\mathcal{L}_M)$.
 - $\varphi \in \text{Cn}(\Gamma)$ iff there exist $\varphi^* \in \mathcal{L}_0$ and $\Gamma^* \subseteq \mathcal{L}_0$ such that
 - (i) φ and Γ can be obtained from φ^* and Γ^* by uniform substitution of sentences.
 - (ii) φ^* is a classical consequence of Γ^* .
- $\Gamma \subseteq \mathcal{L}$ is **consistent** iff there exists $\varphi \in \mathcal{L}_M$, such that $\varphi \notin \text{Cn}(\Gamma)$.

Beliefs and epistemic states

- **E** Set of ‘epistemic states’ (more on these shortly).
- **B** Set of ‘belief sets’, subsets of \mathcal{L}_M that have *at least* the following properties, for all $b \in \mathbf{B}$:
 - **Closure (Cl)** $\text{Cn}(b) \subseteq b$.
 - **Consistency (Con)** b is consistent.
- **Bel** Belief function from **E** to **B**.
- Interpretation: gives us the beliefs that an agent is permitted to hold, given his or her epistemic state.

Orders

- A **preorder** \geq on a set S is a binary relation on S that is both reflexive and transitive.
- \sim The symmetric part of a preorder \geq .
- **max**(S, \geq) The set of maximal elements of S according to \geq , i.e. $\{x \in S : \forall x^* \in S, x \geq x^*\}$.
- W_{n+1} Set of all total preorders over W_n , where $n \in \mathbb{N}_0$.
- **W** Union of the W_i .

Levi's suggestion

- A pair of proposals regarding \diamond (Levi [1988]):
 - ($\diamond 1$) For all $\varphi \in \mathcal{L}_M$ and $x \in \mathbf{E}$, $\neg \diamond \neg \varphi \in \text{Bel}(x)$ iff $\varphi \in \text{Bel}(x)$.
 - ($\diamond 2$) For all $\varphi \in \mathcal{L}_M$ and $x \in \mathbf{E}$, if $\varphi \notin \text{Bel}(x)$, then $\diamond \neg \varphi \in \text{Bel}(x)$.
- ($\diamond 1$) seems clearly correct.
- Note in passing that, conveniently:

Observation 1: Given (Cl) and (Con), ($\diamond 1 \Leftrightarrow$) entails that $\varphi \wedge \diamond \neg \varphi \notin \text{Bel}(x)$, for all $\varphi \in \mathcal{L}_M$ and $x \in \mathbf{E}$. [Proof \triangleright]
- ($\diamond 2$), however, may be more problematic.

Two consequences of ($\diamond 2$)

- Trivially:

Observation 2: ($\diamond 1 \Leftrightarrow$) and ($\diamond 2$) jointly entail (OM). [Proof \triangleright]

Where:

Opinionation wrt Modals (OM) For all $\varphi \in \mathcal{L}_M$ and $x \in \mathbf{E}$, either $\diamond \varphi \in \text{Bel}(x)$ or $\neg \diamond \varphi \in \text{Bel}(x)$.
- Relatedly:

Observation 3: Given (Con), ($\diamond 1 \Leftrightarrow$) and ($\diamond 2$) jointly entail (Red). [Proof \triangleright]

Where:

Reduction (Red) For all $\varphi \in \mathcal{L}_M$ and $x \in \mathbf{E}$:

 - (i) If $\diamond \diamond \varphi \in \text{Bel}(x)$ then $\diamond \varphi \in \text{Bel}(x)$.
 - (ii) If $\diamond \square \varphi \in \text{Bel}(x)$ then $\square \varphi \in \text{Bel}(x)$.

Sorensen's objections

- Sorensen [2009]: (OM) and (Red) seem too strong.
- Regarding (Red):
 - B*: There might be a possibility of still getting that grant.
 - A*: There *is* a possibility that we'll still get the grant?
 - B*: That's not what I said: there *might* be such a possibility...
- Regarding (OM):
 - A*: Do you think that there's a possibility that we will get that grant?
 - B*: I don't know. Perhaps it's already too late.

Introducing ($\diamond 3$)

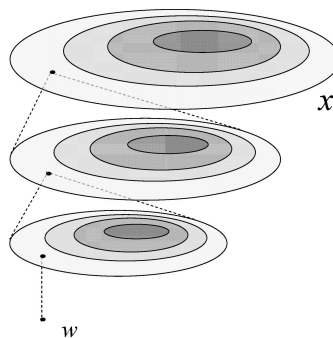
- If we grant that ($\diamond 2$) must go, we are quickly led to the following mild strengthening of its negation (argument omitted):
 - If it is permissible to suspend judgment on φ , then it is optional to do so without accepting $\diamond \neg \varphi$.
- More formally:
 - ($\diamond 3$) For all $\varphi \in \mathcal{L}_M$, there exists $x \in \mathbf{E}$ such that $\varphi, \neg \varphi, \diamond \varphi \notin \text{Bel}(x)$ iff there exists $y \in \mathbf{E}$ such that $\varphi, \neg \varphi \notin \text{Bel}(y)$ and $\diamond \varphi \in \text{Bel}(y)$.
- Question: If ($\diamond 3$) is correct, what impact, if any, does this have on the standard view of epistemic states?

The standard view

- The view in question:
Total Preorder in W_1 (TP₁) $\mathbf{E} = W_1$.
- Equally standardly:
(Fac) For all $\varphi \in \mathcal{L}_0$ and $x \in \mathbf{E}$, $\varphi \in \text{Bel}(x)$ iff $\max(W_0, x) \subseteq \llbracket \varphi \rrbracket$.
- It is easy to establish, however, that:
Observation 4: (TP₁), (Fac) and ($\diamond 3$) are jointly inconsistent.
 [Proof ▷]
- To keep ($\diamond 3$) we'll need to enlarge \mathbf{E} by weakening (TP₁).

Meta-orderings

- One straightforward move that does the job:
Total Preorder in $W - W_0$ (TP) $\mathbf{E} = W - W_0$.
- Illustration, where $x \in W_3$:



Defining Bel

- We then define Bel inductively.
- Basis step:
 - (BS) For all x in W_1 :
 - (a) For all $\varphi \in \mathcal{L}_0$, $\varphi \in \text{Bel}(x)$ if, for all $y \in \max(W_0, x)$, $y \in \llbracket \varphi \rrbracket$.
 - (b) For all $\varphi \in \mathcal{L}_M$, $\Box\varphi \in \text{Bel}(x)$, if $\varphi \in \text{Bel}(x)$.
 - (c) For all $\varphi \in \mathcal{L}_M$, $\Diamond\varphi \in \text{Bel}(x)$ if $\varphi \in \text{Cn}(\{\psi, \chi\})$ for (i) some $\chi \in \text{Bel}(x)$ and (ii) some $\psi \in \mathcal{L}_0$ such that, for some $y \in \max(W_0, x)$, $y \in \llbracket \psi \rrbracket$.
 - (d) For all $\Gamma \subseteq \mathcal{L}_M$, $\text{Cn}(\Gamma) \subseteq \text{Bel}(x)$ if $\Gamma \subseteq \text{Bel}(x)$.
 - (e) For all $\varphi \in \mathcal{L}_M$, $\varphi \in \text{Bel}(x)$ only if its membership can be derived from (a)-(d).
- It can be proven that:

Observation 5: Given (TP₁), (BS) entails (\Diamond 1) and (\Diamond 2). [Proof \triangleright]

Defining Bel (ctd.)

- (\cdot) For all $x \in \mathbf{E}$ and $\varphi \in \mathcal{L}_M$, $x \in \langle \varphi \rangle$ iff $\varphi \in \text{Bel}(x)$.
- Inductive step:
 - (IS) For all x in W_{n+1} ($n \geq 1$):
 - (a) For all $\varphi \in \mathcal{L}_M$, $\varphi \in \text{Bel}(x)$ if, for all $y \in \max(W_n, x)$, $y \in \langle \varphi \rangle$.
 - (b) For all $\varphi \in \mathcal{L}_M$, $\Box\varphi \in \text{Bel}(x)$, if $\varphi \in \text{Bel}(x)$.
 - (c) For all $\varphi \in \mathcal{L}_M$, $\Diamond\varphi \in \text{Bel}(x)$ if $\varphi \in \text{Cn}(\{\psi, \chi\})$ for (i) some $\chi \in \text{Bel}(x)$ and (ii) some $\psi \in \mathcal{L}_M$ such that, for some $y \in \max(W_n, x)$, $y \in \langle \psi \rangle$.
 - (d) For all $\Gamma \subseteq \mathcal{L}_M$, $\text{Cn}(\Gamma) \subseteq \text{Bel}(x)$ if $\Gamma \subseteq \text{Bel}(x)$.
 - (e) For all $\varphi \in \mathcal{L}_M$, $\varphi \in \text{Bel}(x)$ only if its membership can be derived from (a)-(d).
- As promised:

Observation 6: Given (TP), (BS) and (IS) jointly entail (\Diamond 1) and (\Diamond 3). [Proof \triangleright]

Logics

- We can use the two models presented here to define a ‘consequence’ relation on $\wp(\mathcal{L}_M) \times \mathcal{L}_M$.
- \models_M Where $\Gamma \subseteq \mathcal{L}_M$ and $\varphi \in \mathcal{L}_M$, $\Gamma \models_M \varphi$ iff $\varphi \in \text{Bel}(x)$ for all $x \in \mathbf{E}$, such that $\Gamma \subseteq \text{Bel}(x)$.
- \models_M looks very much like supervaluationist global consequence:
 - Observation 7:** Even given (TP_1) , \models_M fails to satisfy (i) contraposition, (ii) conditional proof and (iii) reasoning by cases. [Proof ▷]
- Furthermore:
 - Observation 8:** Given (TP_1) , the S5 axioms are \models_M -valid. [Proof ▷]
- Question: What happens if we retreat to (TP) ?

Concluding comments

- (TP) yields a very large set of epistemic states.
- Modal agnosticism can be similarly accommodated in models that are more quantitatively parsimonious (e.g. epist. states as sets of sets... of elements of W_1).
- However:
 - (a) Such models are arguably not as *qualitatively* parsimonious (orderings + sets vs orderings all the way up)
 - (b) (TP) turns out to have some interesting applications to the issue of *left-nested conditionals*.
- But (b) is another talk altogether...

Thank you!

Questions and comments welcome:
jacob.chandler@hiw.kuleuven.be

References

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Proof of Observation 1

- (1) For some $\varphi \in \mathcal{L}_M$ and $x \in \mathbf{E}$, $\varphi \wedge \Diamond \neg \varphi \in \text{Bel}(x)$ [for reductio]
- (2) $\Diamond \neg \varphi \in \text{Bel}(x)$. [(1), (C1)]
- (3) If $\varphi \in \text{Bel}(x)$, then $\Diamond \neg \varphi \notin \text{Bel}(x)$ [$(\Diamond 1 \Leftarrow)$, (Con)]
- (4) $\varphi \notin \text{Bel}(x)$ [(2), (3)]
- (5) $\varphi \in \text{Bel}(x)$ [(1), (C1)] □

[Back ◀]

Proof of Observation 2

- (1) For all $\varphi \in \mathcal{L}_M$ and $x \in \mathbf{E}$, either (i) $\neg \varphi \in \text{Bel}(x)$ or (ii) $\neg \varphi \notin \text{Bel}(x)$
- (2) Assume (i)
- (3) $\neg \Diamond \varphi \in \text{Bel}(x)$ [$(\Diamond 1 \Leftarrow)$]
- (4) Assume (ii)
- (5) $\Diamond \varphi \in \text{Bel}(x)$ [$(\Diamond 2)$] □

[Back ◀]

Proof of Observation 3

- (1) For some $\varphi \in \mathcal{L}_M$ and $x \in \mathbf{E}$, $\diamond \diamond \varphi \in \text{Bel}(x)$ [for CP]
- (2) $\neg \diamond \neg \neg \diamond \varphi \notin \text{Bel}(x)$ [(1), (Cl), (Con)]
- (3) $\neg \diamond \varphi \notin \text{Bel}(x)$ [(2), contrapositive of $(\diamond 1 \Leftarrow)$]
- (4) $\diamond \varphi \in \text{Bel}(x)$ [(3), $(\diamond 1 \Leftarrow)$ and $(\diamond 2)$ via (OM)] □

[Back ◀]

Proof of Observation 4

- (1) Let $\mathcal{A} = \{\varphi\}$, with $w, w^* \in W_0$, such that $w \models \varphi$ and $w^* \models \neg \varphi$
- (2) There exists a *unique* $x \in \mathbf{E}$ such that $\varphi, \neg \varphi \notin \text{Bel}(x)$: the $x \in \mathbf{E}$, such that $w \sim_x w^*$ [(1), (TP₁), (Fac)]
- (3) It isn't the case that $\diamond \varphi \in \text{Bel}(x)$ and $\diamond \varphi \notin \text{Bel}(x)$
- (4) $(\diamond 3)$ is false [(2), (3)] □

[Back ◀]

Proof of Observation 5

($\diamond 1$): Given (TP₁), (BS)(b) is the only way to secure membership of $\text{Bel}(x)$ for any $\Box\varphi$, such that $\varphi \in \mathcal{L}_M$, for any $x \in \mathbf{E}$. The desired conclusion then follows from (BS)(e).

($\diamond 2$): We first define a function $d : \mathcal{L}_M \mapsto \mathbb{N}_0$ as follows:

- (i) For all $\varphi \in \mathcal{L}_0$, $d(\varphi) = 0$
- (ii) For all $\varphi \in \mathcal{L}_M$, $d(\neg\varphi) = d(\varphi)$
- (iii) For all $\varphi, \psi \in \mathcal{L}_M$, $d(\varphi \vee \psi) = d(\varphi \wedge \psi) = d(\varphi \rightarrow \psi) = \max\{d(\varphi), d(\psi)\}$
- (iv) For all $\varphi \in \mathcal{L}_M$, $d(\diamond\varphi) = d(\varphi) + 1$

We now define:

$$\mathcal{L}_n := \{\varphi \in \mathcal{L}_M : d(\varphi) \leq n\}$$

Note that $\mathcal{L}_M = \bigcup \mathcal{L}_n$, $n \in \mathbb{N}_0$.

Proof of Observation 5 (ctd.)

We also define the corresponding restricted version of ($\diamond 2$):

$$(\mathcal{L}_n \diamond 2) \text{ For all } \varphi \in \mathcal{L}_n \text{ and all } x \in \mathbf{E}, \text{ if } \varphi \notin \text{Bel}(x), \text{ then } \diamond\neg\varphi \in \text{Bel}(x).$$

We first prove ($\mathcal{L}_0 \diamond 2$). Assume for CP that $\varphi \notin \text{Bel}(x)$, where $\varphi \in \mathcal{L}_0$. It follows, given the contrapositive of (BS)(a), that, for some $y \in \max(W_0, x)$, $y \in \llbracket \neg\varphi \rrbracket$. From this, given (BS)(c), we recover the fact that $\diamond\neg\varphi \in \text{Bel}(x)$.

We now prove that if ($\mathcal{L}_n \diamond 2$), then ($\mathcal{L}_{n+1} \diamond 2$), where $n \in \mathbb{N}_0$.

Consider an arbitrary $\varphi \in \mathcal{L}_{n+1}$. Let $\text{DNF}(\varphi)$ denote its DNF, i.e. its equivalent disjunction of conjunctions of sentences ψ_i , such that (i) ψ_i is a literal or (ii) $\psi_i = \diamond\chi$ or $\psi_i = \neg\diamond\chi$, where $\chi \in \mathcal{L}_n$:

$$\text{DNF}(\varphi) = (\psi_1, \wedge \dots) \vee (\psi_n, \wedge \dots) \vee \dots$$

Proof of Observation 5 (ctd.)

Call a sentence ‘indefinite’ iff neither it nor its negation is in $\text{Bel}(x)$.

$\varphi \notin \text{Bel}(x)$ iff either (a) φ is definite and $\neg\varphi \in \text{Bel}(x)$ or (b) φ is indefinite.

If (a), it immediately follows, by (IS)(c), that $\diamond\neg\varphi \in \text{Bel}(x)$.

Assume $(\mathcal{L}_n \diamond 2)$. It follows from that, alongside $(\diamond 1)$, that, for all $\chi \in \mathcal{L}_n$ and $x \in \mathbf{E}$, either $\diamond\chi \in \text{Bel}(x)$ or $\neg\diamond\chi \in \text{Bel}(x)$ (see proof of Obs 2).

In other words: all non-literal conjuncts in $\text{DNF}(\varphi)$ are definite.

Proof of Observation 5 (ctd.)

So if (b), then $\text{DNF}(\varphi)$ must (i) include at least one indefinite disjunct, which itself, by the previous remark, must contain at least one indefinite *literal* conjunct, and (ii) only include definite disjuncts whose negation is in $\text{Bel}(x)$.

Now assume (b) and consider $\neg\text{DNF}(\varphi)$. This will be equivalent to a conjunction α of disjunctions that either (i) contain an indefinite literal disjunct, namely the negation of the corresponding conjunct in $\text{DNF}(\varphi)$, or (ii) are members of $\text{Bel}(x)$:

$$\alpha = (\neg\psi_1, \vee \dots) \wedge (\neg\psi_n, \vee \dots) \wedge \dots$$

Proof of Observation 5 (ctd.)

Let Γ denote the set of indefinite literals in α . Since $\text{DNF}(\varphi) \notin \text{Bel}(x)$, it follows that the disjunction of their negations isn't in $\text{Bel}(x)$. It then follows that there exists $y \in \max(W_0, x)$ such that $y \in \llbracket \wedge \Gamma \rrbracket$.

Now it is easy to show that α is a joint consequence of $\wedge \Gamma$ and the conjunction of the conjuncts in α that are members of $\text{Bel}(x)$ (if any).

It then follows from (BS)(c) that $\diamond \alpha \in \text{Bel}(x)$ and hence, by (BS)(d), that $\diamond \neg \varphi \in \text{Bel}(x)$.

We can therefore conclude that $(\mathcal{L}_{n+1} \diamond 2)$ holds. \square

[Back \triangleleft]

Proof of Observation 6

($\diamond 1$): Call $(W_n \diamond 1)$, the restriction of ($\diamond 1$) to W_n . We saw in Observation 5 that (BS) entails that $(W_1 \diamond 1)$ holds, given (BS)(e) and the fact that (BS)(b) is the only way to secure membership of $\Box \varphi$ when $x \in W_1$. For similar reasons, (IS) entails that $(W_n \diamond 1)$, for $n \geq 2$.

($\diamond 3 \Leftarrow$): We prove this for $x \in W_n$, where $n \geq 2$; the proof for $n = 1$ is analogous (simply swap $\llbracket \cdot \rrbracket$ for $\langle \cdot \rangle$).

Assume for CP that there exists $y \in W_n$ such that $\varphi, \neg \varphi \notin \text{Bel}(y)$ and $\diamond \varphi \in \text{Bel}(y)$.

It follows from (IS)(a) and the fact that $\varphi \notin \text{Bel}(y)$, that there exists $x \in W_{n-1}$ such that $x \in \max(W_{n-1})$ and $x \notin \langle \varphi \rangle$.

Proof of Observation 6 (ctd.)

There exists $y^* \in W_n$ such that $\max(W_{n-1}, y^*) = \{x\}$, as well as $z \in W_{n+1}$ such that $\max(W_n, z) = \{y, y^*\}$.

Since, as is easily verified, $\diamond\varphi \notin \text{Bel}(y^*)$ and since $\varphi, \neg\varphi \notin \text{Bel}(y)$, it follows that $\varphi, \neg\varphi, \diamond\varphi \notin \text{Bel}(z)$.

($\diamond 3 \Rightarrow$): Assume for CP that there exists $y \in W_n$ such that $\varphi, \neg\varphi, \diamond\varphi \notin \text{Bel}(y)$.

From the fact that $\neg\varphi \notin \text{Bel}(y)$, it follows that, for some $x \in W_1$, $\neg\varphi \notin \text{Bel}(x)$. Indeed, assume for reductio, that there is no such x . It then follows that $\neg\varphi \in \text{Bel}(y)$, contrary to our initial assumption, since if for all $x \in W_n$, $\neg\varphi \in \text{Bel}(x)$, then, trivially, for all $x^* \in W_{n+1}$, for all $x \in \max(W_n, x^*)$, $x \in \langle \neg\varphi \rangle$ and hence by (IS)(a), $\neg\varphi \in \text{Bel}(x^*)$.

Proof of Observation 6 (ctd.)

Either (i) $\varphi \in \text{Bel}(x)$ or (ii) $\varphi \notin \text{Bel}(x)$.

Assume (i). Now:

Observation 9: If there exists $x \in W_n$ ($n \geq 1$) such that $\varphi \in \text{Bel}(x)$, then there exists $x^* \in W_{n+1}$ such that $\varphi \in \text{Bel}(x^*)$.

Indeed, consider any x^* such that $\max(W_n, x^*) = \{x\}$.

So, by Obs 9, there exists $y^* \in W_n$ such that $\varphi \in \text{Bel}(y^*)$. There also exists $z \in W_{n+1}$, such that $\max(W_n, z) = \{y, y^*\}$.

Since $\varphi \in \text{Bel}(y^*)$, by (IS)(c), $\diamond\varphi \in \text{Bel}(z)$. Since $\varphi, \neg\varphi \notin \text{Bel}(y)$, by (IS), $\varphi, \neg\varphi \notin \text{Bel}(z)$.

Assume (ii). By ($\diamond 2$), which holds for all $x \in W_1$ (see Observation 5), since $\neg\varphi \notin \text{Bel}(x)$, $\diamond\varphi \in \text{Bel}(x)$. □

Proof of Observation 7

- (i) For all $\varphi \in \mathcal{L}_M$, $\varphi \models_M \Box\varphi$. However, it is not the case that for all $\varphi \in \mathcal{L}_M$, $\Diamond\neg\varphi \models_M \neg\varphi$. Countermodel: see epistemic state x in proof of Observation 4.
- (ii) For all $\varphi \in \mathcal{L}_M$, $\varphi \models_M \Box\varphi$. However, it is not the case that for all $\varphi \in \mathcal{L}_M$, $\models_M \varphi \rightarrow \Box\varphi$. Countermodel: same as above.
- (iii) For all $\varphi \in \mathcal{L}_M$, $\varphi \models_M \Box\varphi \vee \Box\neg\varphi$ and $\neg\varphi \models_M \Box\varphi \vee \Box\neg\varphi$. However, it is not the case that for all $\varphi \in \mathcal{L}_M$, $\varphi \vee \neg\varphi \models_M \Box\varphi \vee \Box\neg\varphi$. Countermodel: same as above. \square

[Back \triangleleft]

Proof of Observation 8

$\models_M \mathbf{K}$: Either (i) $\Diamond\neg\varphi \in \text{Bel}(x)$ or (ii) $\Diamond\neg\varphi \notin \text{Bel}(x)$. Furthermore, either (iii) $\neg\Diamond\neg\psi \in \text{Bel}(x)$ or (iv) $\neg\Diamond\neg\psi \notin \text{Bel}(x)$.

Now from (TP₁) and (BS), we recover (\Diamond 1) and (\Diamond 2) (see Obs. 5). From these, we obtain (OM) (see Obs. 2).

If (iv) then, by (OM), $\Diamond\neg\psi \in \text{Bel}(x)$. If, furthermore, (ii) then, again by (OM), it follows that $\neg\Diamond\neg\varphi \in \text{Bel}(x)$ and therefore, by (BS), that $\varphi \in \text{Bel}(x)$. Finally, by (BS)(c), $\Diamond(\varphi \wedge \neg\psi) \in \text{Bel}(x)$.

If (i), or (ii), then $\Diamond\neg\varphi \vee \neg\Diamond\neg\psi \in \text{Bel}(x)$

So, whatever way, $\Diamond(\varphi \wedge \neg\psi) \vee (\Diamond\neg\varphi \vee \neg\Diamond\neg\psi) \in \text{Bel}(x)$, and hence, by (BS)(d), $\Box(\varphi \rightarrow \psi) \rightarrow (\Box\varphi \rightarrow \Box\psi) \in \text{Bel}(x)$.

Proof of Observation 8 (ctd.)

$\models_M \mathbf{T}$: By $(\diamond 2)$, either $\varphi \in \text{Bel}(x)$ or $\diamond \neg \varphi \in \text{Bel}(x)$ and hence, by (BS)(d) $\Box \varphi \rightarrow \varphi \in \text{Bel}(x)$.

$\models_M \mathbf{5}$: By (OM), either $\diamond \varphi \in \text{Bel}(x)$ or $\neg \diamond \varphi \in \text{Bel}(x)$. By (BS)(b) and (IS)(b), if $\diamond \varphi \in \text{Bel}(x)$, then $\Box \diamond \varphi \in \text{Bel}(x)$. Therefore, by (BS)(d), $\diamond \varphi \rightarrow \Box \diamond \varphi \in \text{Bel}(x)$. \square

[Back \triangleleft]